

② Unit axiom: the functor of left and right mult.
by 1:
 $L_1: X \mapsto 1 \otimes X$, $R_1: X \mapsto X \otimes 1$
are autoequivalences on $\mathcal{C}$.

Lemma: There are natural isomorphisms

$$l: L_1 \rightarrow \text{Id}_{\mathcal{C}} , \quad r: R_1 \rightarrow \text{Id}_{\mathcal{C}}.$$

proof: We will prove that for l (r is similar).

We will make such an ℓ such that:

$$\begin{array}{ccc}
 1 \otimes (1 \otimes X) & \xrightarrow{\bar{a}'_{1,1,X}} & (1 \otimes 1) \otimes X \\
 \searrow L_1(l_X) & & \swarrow L \otimes \text{id}_X \\
 & 1 \otimes X &
 \end{array}
 \quad L: 1 \otimes 1 \rightarrow 1$$

We want $L_1(l_X) = (L \otimes \text{id}_X) \circ \bar{a}'_{1,1,X}$.

But, as we proved yesterday, because L_1 is an equivalence of category, it induces a bijection on the hom-set, in particular, $l_X = L_1^{-1}[(L \otimes \text{id}_X) \circ \bar{a}'_{1,1,X}]$.

$\rightarrow l_X$ is isom. OK. inverse in $\uparrow$ hom-set

Such an l defined by above l_x is a natural transformation if by def: $\forall X \xrightarrow{f} Y$

$$\begin{array}{ccc} L_1(X) & \xrightarrow{l_x} & Id_Y(X) \\ L_1(f) \downarrow & & \downarrow Id_Y(f) \end{array} \text{ commutes.}$$

$$L_1(Y) \xrightarrow{l_y} Id_Y(Y)$$

$$\Leftrightarrow \begin{array}{ccc} L_1 \circ L_1(X) & \xrightarrow{L_1(l_x)} & L_1 \circ Id_Y(X) \\ L_1 \circ L_1(f) \downarrow & & \downarrow L_1 \circ Id_Y(f) \end{array} \text{ commutes}$$

$$L_1 \circ L_1(Y) \xrightarrow{L_1(l_y)} L_1 \circ Id_Y(Y)$$

the last diagram reformulates and decomposes as follows:

$$\begin{array}{ccccc}
 1 \otimes (1 \otimes X) & \xrightarrow{\bar{a}_{1,1,1,X}'} & (1 \otimes 1) \otimes X & \xrightarrow{\iota \otimes \text{id}_X} & 1 \otimes X \\
 \downarrow \text{Id} \otimes (\text{Id} \otimes f) & & \downarrow (\text{id} \otimes \text{id}) \otimes f & & \downarrow \text{Id} \otimes f \\
 1 \otimes (1 \otimes Y) & \xrightarrow{\bar{a}_{1,1,Y}'} & (1 \otimes 1) \otimes Y & \xrightarrow{\iota \otimes \text{id}_Y} & 1 \otimes Y
 \end{array}$$

But the left and right squares commute because

left: $\bar{a}_{1,1,1,-}'$ is a natural transformation (OK).

right: $(\text{id} \otimes f) \circ (\iota \otimes \text{id}_X) = \iota \otimes f$. $\left\{ \begin{array}{l} \text{The result} \\ \text{follows.} \end{array} \right.$

$(\iota \otimes \text{id}_Y) \circ (\underbrace{\text{id} \otimes \text{id}_1}_{\text{id}_1}) \otimes f = \iota \otimes f$

Exo:

$$\overline{\text{Id}_{X \otimes Y}} = \overline{\text{Id}_X \otimes \text{Id}_Y}$$

Proposition: $\forall X \in \mathcal{C}$

$$l_{1 \otimes X} = \text{id}_1 \otimes l_X \quad , \quad r_{X \otimes 1} = r_X \otimes \text{id}_1$$

$\left\{ \begin{array}{l} \text{!} \\ l_{X \otimes Y} \neq l_X \otimes l_Y \\ \text{in general.} \end{array} \right.$

Proof: We know that l is natural transformation, i.e. $l: 1 \rightarrow \text{id}$

$L_1(X)$	$\xrightarrow{l_X}$	$\text{Id}_1(X)$	} Apply it to $f = l_Z$ ($X = 1 \otimes Z, Y = Z$)
$f \downarrow$		$\downarrow \text{Id}_1(f)$	
$L_1(Y)$	$\xrightarrow{l_Y}$	$\text{Id}_1(Y)$	
commutes.			

$1 \otimes (1 \otimes Z)$	$\xrightarrow{l_{1 \otimes Z}}$	$1 \otimes Z$	} $\Leftrightarrow l_Z \circ l_{1 \otimes Z}$
$\downarrow \text{id} \otimes l_Z$		$\downarrow l_Z$	
$1 \otimes Z$	$\xrightarrow{l_Z}$	Z	
commutes			

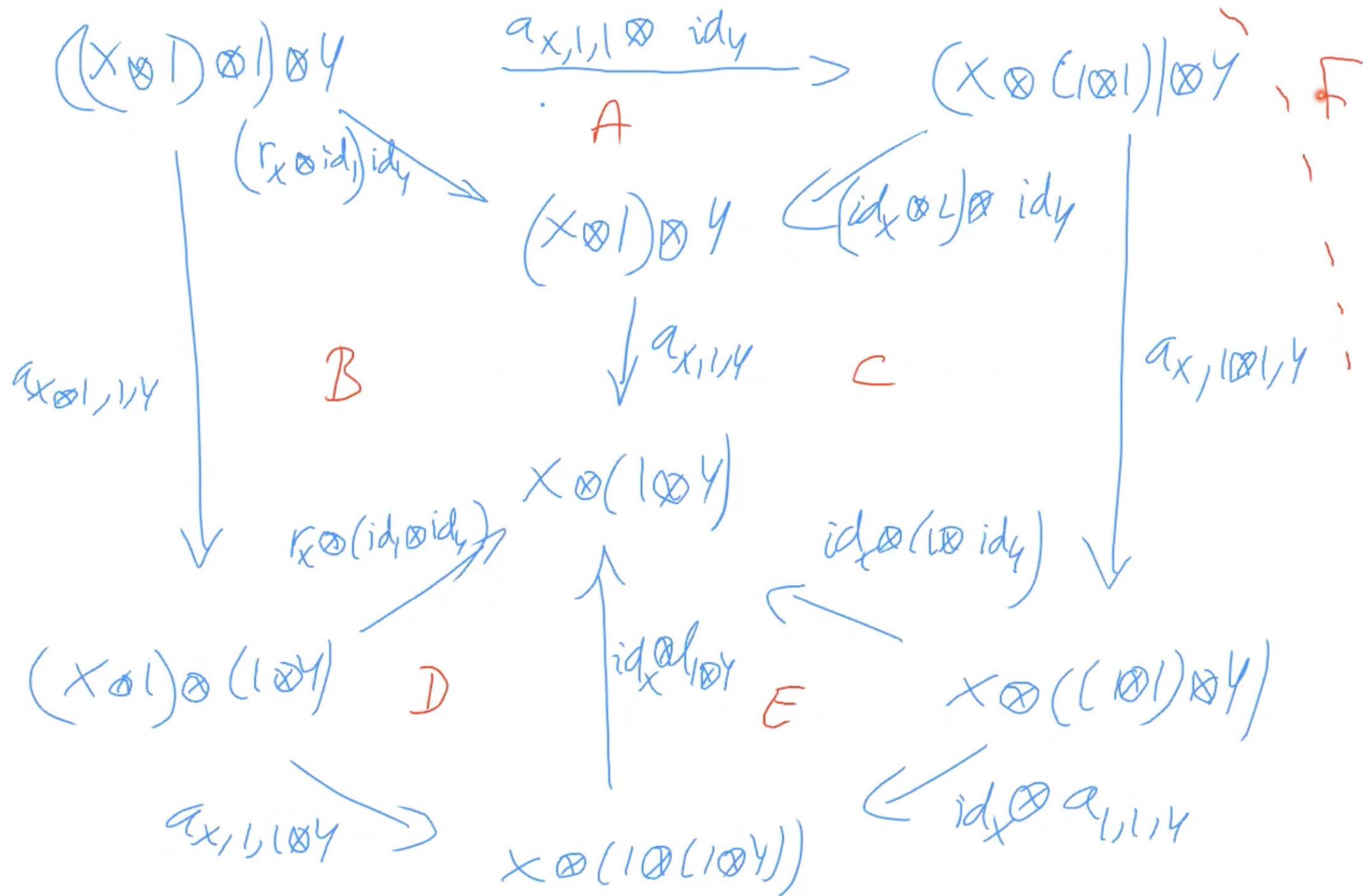
$l_Z \circ (id_1 \otimes l_Z)$	} but l_Z is also $\left\{ \begin{array}{l} \text{you can} \\ \text{apply } l_Z \end{array} \right.$
So the expected equality holds $\square$	

Proposition (triangle diagram) : $\forall X, Y \in \mathcal{C}$

$$\begin{array}{ccc} (X \otimes 1) \otimes Y & \xrightarrow{a_{X,1,Y}} & X \otimes (1 \otimes Y) \\ & \searrow \scriptstyle X \otimes \text{id}_Y & \swarrow \scriptstyle \text{id}_X \otimes Y \\ & X \otimes Y & \end{array}$$

commutes.

Rh: The former definition of a monoidal category is Pentagon axiom + triangle axiom. But the definition used in this course is Pent. axiom + Unit axiom (and above prove the triangle axiom)



Ⓕ: is just the pentagon axiom.

ⒼⓈ: just apply that α is a natural transformation.

Ⓐ: Need to have $r_x \otimes \text{id}_l \stackrel{?}{=} (\text{id}_x \otimes \iota) \circ \alpha_{x,1,l}$

But: $R_1(r_x) = r_{x \otimes 1} = r_x \otimes \text{id}_l$ (by previous proposition)

and: $R_1(r_x) = (\text{id}_x \otimes \iota) \circ \alpha_{x,1,l}$ by definition Ⓢ.

Ⓔ: As for Ⓐ but for l . Ⓢ.

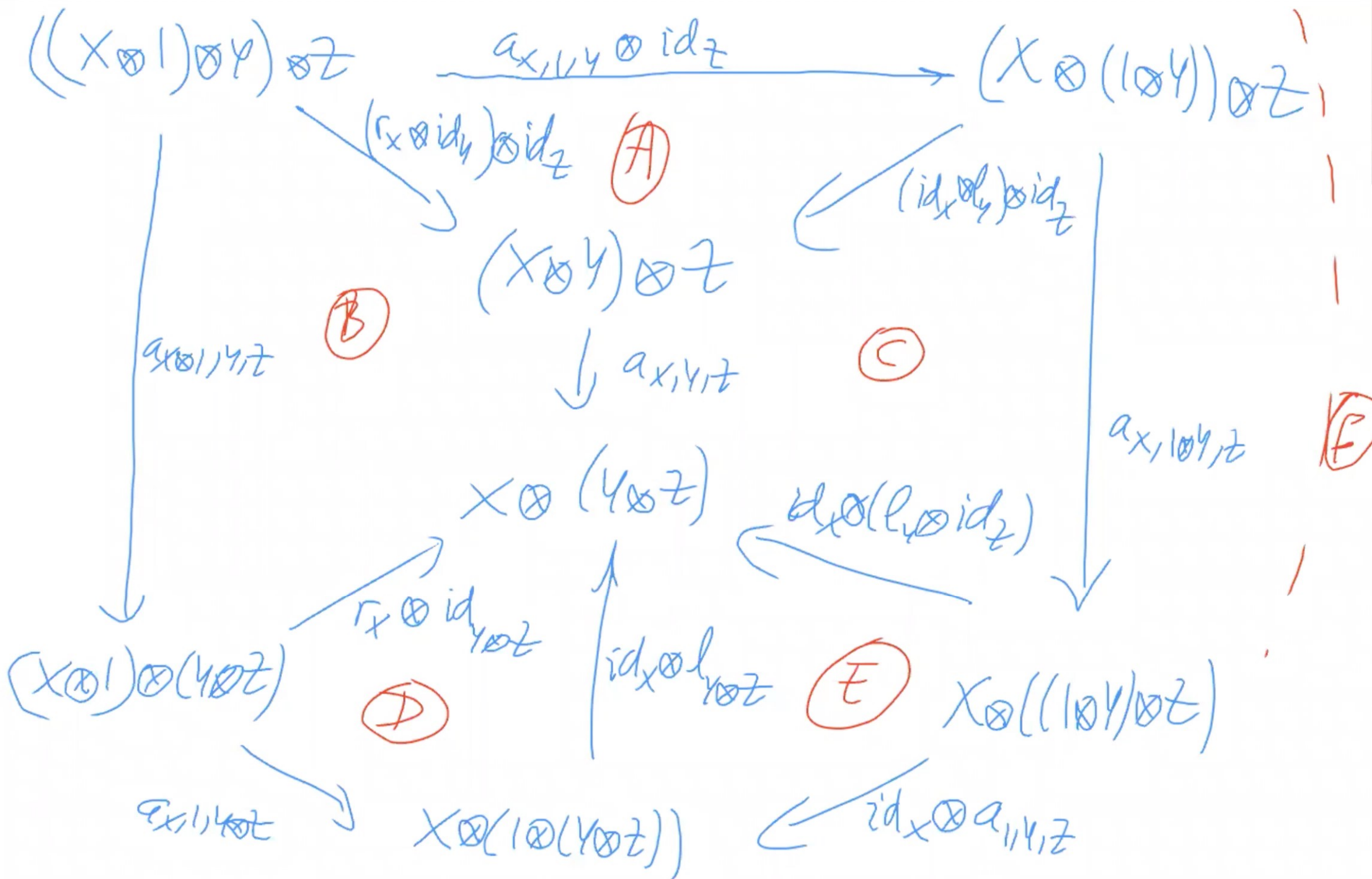
⇒ Ⓓ holds, so the result follows because L_1 is an equivalence of ab. $\square$

Proposition: $\forall X, Y \in \mathcal{V}$, the following two diagrams commute:

$$\begin{array}{ccc}
 (1 \otimes X) \otimes Y & \xrightarrow{a_{1, X, Y}} & 1 \otimes (X \otimes Y) \\
 \downarrow l_X \otimes \text{id}_Y & & \swarrow l_{X \otimes Y} \\
 & X \otimes Y &
 \end{array}$$

$$\begin{array}{ccc}
 (X \otimes Y) \otimes 1 & \xrightarrow{a_{X, Y, 1}} & X \otimes (Y \otimes 1) \\
 \downarrow r_{X \otimes Y} & & \swarrow \text{id}_X \otimes r_Y \\
 & X \otimes Y &
 \end{array}$$

Proof: Let us prove the first one (the second is similar).



As before, (F): pentagon axiom.

(B), (C): a is a natural transformation $\left(\begin{array}{c} \text{with} \\ \text{id}_Y \circ z \\ = \\ \text{id}_Y \otimes \text{id}_Z \end{array} \right)$

(A), (D): apply previous proposition (triangle diagram).

So $\textcircled{E}$ commutes.

But $\textcircled{E}$ is what we want after the application of the functor L_1 which is an equivalence of category.

So the result follows. $\square$