

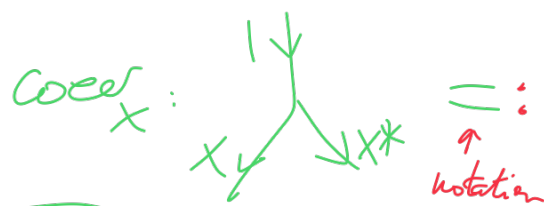
Rigid monoidal categories.

A monoidal category $\mathcal{C}$ is called rigid if every object admits a left and a right dual object.

Def: An object $X^* \in \mathcal{C}$ is called a left dual of $X \in \mathcal{C}$ if there are : $ev_X : X^* \otimes X \rightarrow 1$ (called evaluation morphism)

$coev_X : 1 \rightarrow X \otimes X^*$ (called coevaluation morphism)

Pictures for these morphisms:



such that the following compositions are identity morphisms:

$$\begin{array}{c}
 X \\
 \parallel \\
 1 \otimes X
 \end{array}
 \xrightarrow{\text{coev}_X \otimes \text{id}_X} (X \otimes X^*) \otimes X \xrightarrow{a_{X, X^*, X}} X \otimes (X^* \otimes X) \xrightarrow{\text{id}_X \otimes \text{ev}_X} X = X \otimes 1$$

$$\begin{array}{c}
 X^* \\
 \parallel \\
 X^* \otimes 1
 \end{array}
 \xrightarrow{\text{id}_{X^*} \otimes \text{coev}_X} X^* \otimes (X \otimes X^*) \xrightarrow{a_{X^*, X, X^*}^{-1}} (X^* \otimes X) \otimes X^* \xrightarrow{\text{ev}_X \otimes \text{id}_{X^*}} X^* = 1 \otimes X^*$$

Pictorially, these compositions are as follows: (Zigzag/snake equations)

associativity

$$\text{Diagram 1} = \text{Diagram 2}$$

$$\text{Diagram 3} = \text{Diagram 4}$$

Idem for right dual:

Def: An object *X in $\mathcal{C}$ is called a right dual of X

If there are morphisms:

$$ev'_X : X \otimes {}^*X \rightarrow I$$

$$coev'_X : I \rightarrow {}^*X \otimes X$$

such that

$$X \xrightarrow{id_X \otimes coev'_X} X \otimes ({}^*X \otimes X) \xrightarrow{a^{-1}_{X, {}^*X, X}} (X \otimes {}^*X) \otimes X \xrightarrow{ev'_X \otimes id_X} X$$

$${}^*X \xrightarrow{coev'_X \otimes id_X} ({}^*X \otimes X) \otimes {}^*X \xrightarrow{a_{{}^*X, X, {}^*X}} {}^*X \otimes (X \otimes {}^*X) \xrightarrow{id_{{}^*X} \otimes ev'_X} {}^*X$$

are identity morphisms.

Pictures

$$\begin{array}{c} X \downarrow \quad \swarrow X^* \\ I \downarrow \end{array} =: \begin{array}{c} \text{cup} \\ X \end{array}$$

$$\begin{array}{c} I \downarrow \\ \swarrow X \quad \searrow X \end{array} =: \begin{array}{c} \text{cap} \\ X \end{array}$$

Pictures.

$$\begin{array}{c} X \downarrow \\ \text{coev}'_X \\ \text{cup} \\ \text{ev}'_X \\ X \downarrow \end{array} = \downarrow id_X$$

$$\begin{array}{c} \text{coev}'_X \\ \text{cap} \\ \text{ev}'_X \\ {}^*X \downarrow \end{array} = \downarrow id_{{}^*X}$$

Remark: If X^* is a left dual of X , then X is a right dual of X^* (in other words, you can take $*(X^*) = X$).

Proof: take ${}^*(X^*) := X$, $ev'_{X^*} := ev_X$, $cov'_{X^*} := cov_X$.

Then we need to check that those two compositions are identities

$$\begin{array}{ccccccc}
 X^* & \xrightarrow{\text{id}_{X^*} \otimes \omega_{V'}'} & X^* \otimes ({}^*(X^*) \otimes X^*) & \xrightarrow{\bar{a}_{X^*, {}^*(X^*)}^{-1}} & X^* \otimes (X \otimes ({}^*(X^*) \otimes X^*)) & \xrightarrow{e_{X^*}' \otimes \text{id}_{X^*}} & X^* \\
 X & \xrightarrow{\quad \quad \quad} & \quad \quad \quad & \quad \quad \quad & \quad \quad \quad & \quad \quad \quad & X
 \end{array}$$

They reformulate by the def. as

$$X^* \xrightarrow{\text{id}_{X^*} \otimes \text{coev}_X} X^* \otimes (X \otimes X^*) \xrightarrow{\alpha_{X^*, X, X^*}^{-1}} (X^* \otimes X) \otimes X^* \xrightarrow{\text{ev}_X \otimes \text{id}_{X^*}} X^*$$

which are identities by def of the fact that X^* is a left ideal of X $\square$

Idem: If X^* is the right dual of X , then X is a left dual of X^*
In other words, you can take $(X^*)^* = X$.

Remark: $1^* = {}^*1 = 1$ (by convention $| \otimes | = |$, so just take
ev. and cov. to be identities
(without convention, just take L or $\bar{c}'$)

added part

Left and right adjoint of a functors:

Let $\mathcal{C}$ and $\mathcal{D}$ be two categories (locally small, i.e. $\text{hom}_{\mathcal{C}}(X, Y)$ is a set)

Consider the functors $F: \mathcal{D} \rightarrow \mathcal{C}$, $G: \mathcal{C} \rightarrow \mathcal{D}$

such that for all $X \in \mathcal{C}$, $Y \in \mathcal{D}$

$$\text{hom}_{\mathcal{C}}(F(Y), X) \simeq \text{hom}_{\mathcal{D}}(Y, G(X))$$

Where $\simeq$ is natural in the following sense: there are natural iso. α_X, β_Y

$$\text{hom}_{\mathcal{C}}(F(-), X) : \mathcal{D} \rightarrow \text{Set}$$

$$\Downarrow \alpha_X$$

$$\text{hom}_{\mathcal{D}}(-, G(X)) : \mathcal{D} \rightarrow \text{Set}$$

$$\text{hom}_{\mathcal{C}}(F(Y), -) : \mathcal{C} \rightarrow \text{Set}$$

$$\Downarrow \beta_Y$$

$$\text{hom}_{\mathcal{D}}(Y, G(-)) : \mathcal{C} \rightarrow \text{Set}$$

(and)

It looks like the notion of adjoint matrix/operators defined as
 V : vector space (or Hilbert space) $M \in \text{End}(V)$ matrix.
 $\forall u, v \in V \quad \langle Mu, v \rangle = \langle u, M^*v \rangle$ where M^* = adjoint of M .

F is called a left adjoint functor of G

G ————— right ————— F .

Exercise: Let $\mathcal{C}$ be a category and $\text{End}(\mathcal{C})$ the monoidal category of functors from $\mathcal{C}$ to $\mathcal{C}$ (endofunctors).

Let $F, G \in \text{End}(\mathcal{C})$ be endofunctors.
Show that F is a left adjoint functor of G in $\mathcal{C}$. (def as above with $\mathcal{D} = \mathcal{C}$)

iff F is a left dual of G in $\text{End}(\mathcal{C})$.

Idem: $G \text{ --- right --- } F \text{ ---}$
 $\text{--- } G \text{ --- right --- } F \text{ ---}$

Then: the notion of left (resp. right) dual object generalizes
the notion of adjoint functor.

end of
added part.

Proposition: If $X \in \mathcal{C}$ (monoidal category) has a left (resp. right) dual object, then it is unique up to a unique isomorphism preserving the ev. and coev. morphisms.

proof: Let X_1^*, X_2^* be two left dual objects of X .

Let e_1, c_1, e_2, c_2 be correspondingly evaluation and coev. morphisms.

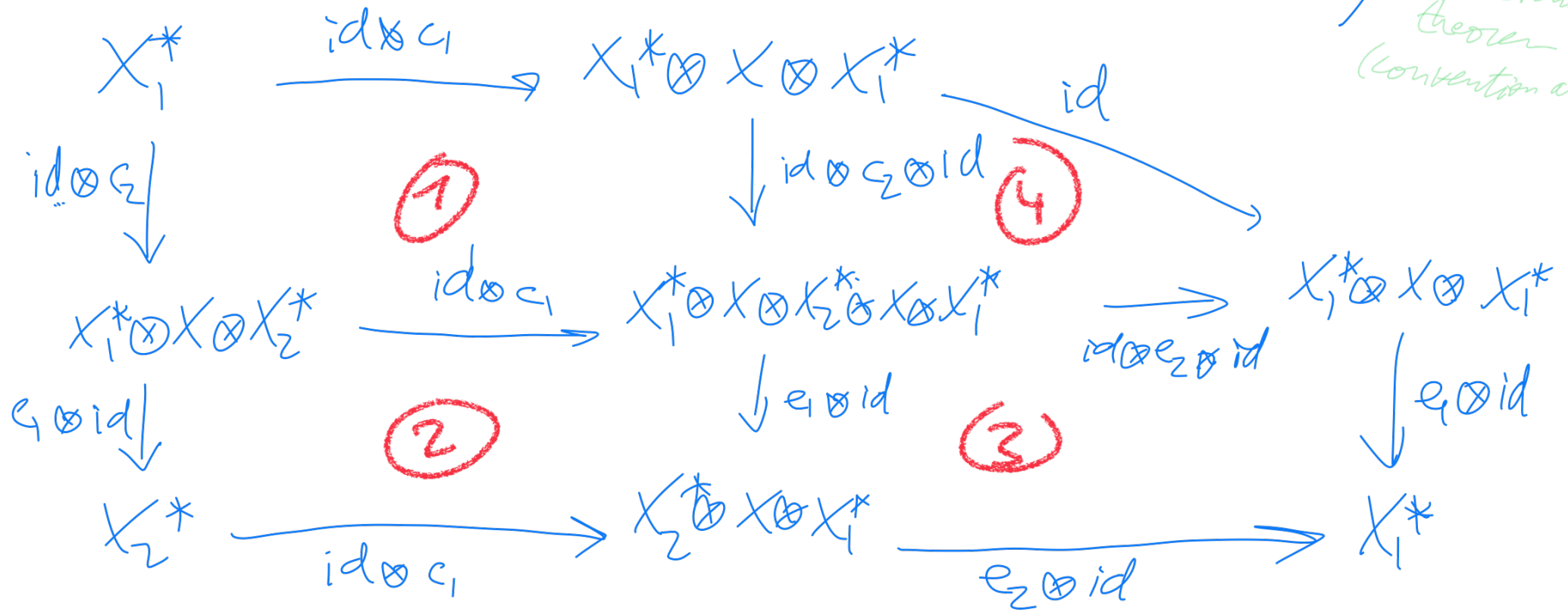
Consider the following morphism $\alpha: X_1^* \rightarrow X_2^*$ defined as

$$\underbrace{X_1^*}_{X_1^* \otimes 1} \xrightarrow{\text{id}_{X_1^*} \otimes c_2} X_1^* \otimes (X \otimes X_2^*) \xrightarrow{\alpha_{X_1^*, X, X_2^*}^{-1}} (X_1^* \otimes X) \otimes X_2^* \xrightarrow{e_1 \otimes \text{id}_{X_2^*}} \underbrace{X_2^*}_{1 \otimes X_2^*}.$$

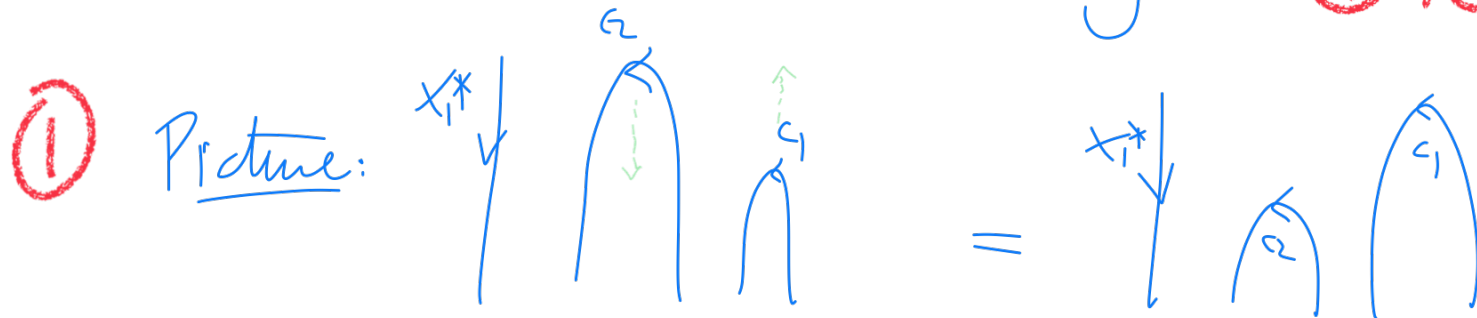
Similarly, consider $\beta: X_2^* \rightarrow X_1^*$.

We will show that $\alpha \circ \beta$ and $\beta \circ \alpha$ are identity morphisms (so that α is iso.)

Consider the following diagram (where the associativity constraints and parenthesize are suppressed) ← OK by coherence theorem (convention also).



Let us show that the local diagrams ①, ②, ③, ④ commute.



(We will see the formal proof tomorrow.)

