

## Lecture 2

16 March 2022

- Goal: Derive Laplace's eq. from electric network.

Given  $u: V \rightarrow \mathbb{R}$ ,  $\Delta u: V \rightarrow \mathbb{R}$  s.t.

$$(\Delta u)_x = \sum_{y \in V(x)} C_{xy} (u_y - u_x)$$

Def

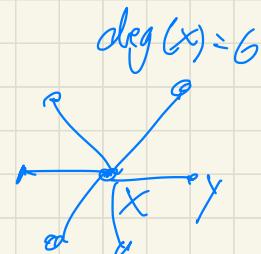
$\Gamma$  graph  $\Gamma = (V, E)$

$V$  vertices,  $E \subset V \times V$  edges (unoriented)

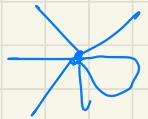
$x, y \in V$   $(x, y) \in E \Leftrightarrow (y, x) \in E$

We say  $x \sim y$  if  $(x, y) \in E$ .

$$N(x) := \{y \in V \mid x \sim y\}$$



$x \in V(x) \Leftrightarrow \exists$  self loop



$\deg(x) = \text{Cardinality of } V(x)$

$\Gamma$  is locally finite if  $\deg(x) < \infty$  for  $x \in V$

Def.

$$\Gamma_1 = (V_1, E_1), \quad \Gamma_2 = (V_2, E_2)$$

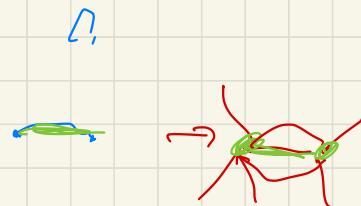
$\phi$  is a map from  $\Gamma_1$  to  $\Gamma_2$  if

$$\phi: V_1 \rightarrow V_2$$

$$E_1 \rightarrow E_2$$

s.t.

$$x \sim y \in V_1 \Rightarrow \phi(x) \sim \phi(y) \in V_2$$



If  $\phi$  has inverse, we say  $\phi$  is isomorphism.

Eg.

$$\mathbb{P} = \mathbb{Z} = \mathbb{Z}'$$



$$\mathbb{P} = \mathbb{Z}^n$$

$\mathbb{Z}_+$



Def (1) An infinite path in  $\mathbb{P}$  is a subgraph isomorphic to  $\mathbb{Z}$ .

(2) One-ended infinite path in  $\mathbb{P}$  is a subgraph isomorphic to  $\mathbb{Z}_+$ .

(3), path of length  $n$  is a subgraph isomorphic to

Def  $\Gamma$  is path connected if

for any two  $x, y \in V$ ,  $\exists$  a path  
connecting  $x$  to  $y$ .

Assumption 1

Remark  $\Gamma$  is assumed to be path connected.

Def Combinatorial distance  $d(x, y)$  is the minimum n s.t.  $\exists$  path of length n connecting to  $x, y \in V$ .

Ass 2.

Remark We assume  $\Gamma$  is countable  
( $\Rightarrow V$  is countable.)

Ex. Locally finite  $\Rightarrow \Gamma$  is countable,

Assumption 3.

Def.,  $(\Gamma, r)$

$r: \mathbb{Y} \rightarrow \mathbb{R}_{>0}$  s.t.  $r(x,y) = r(y,x)$

and

$$C(x) := \sum_{y \in V(x)} \frac{1}{r(x,y)} < \infty \quad \text{for all } x \in V.$$

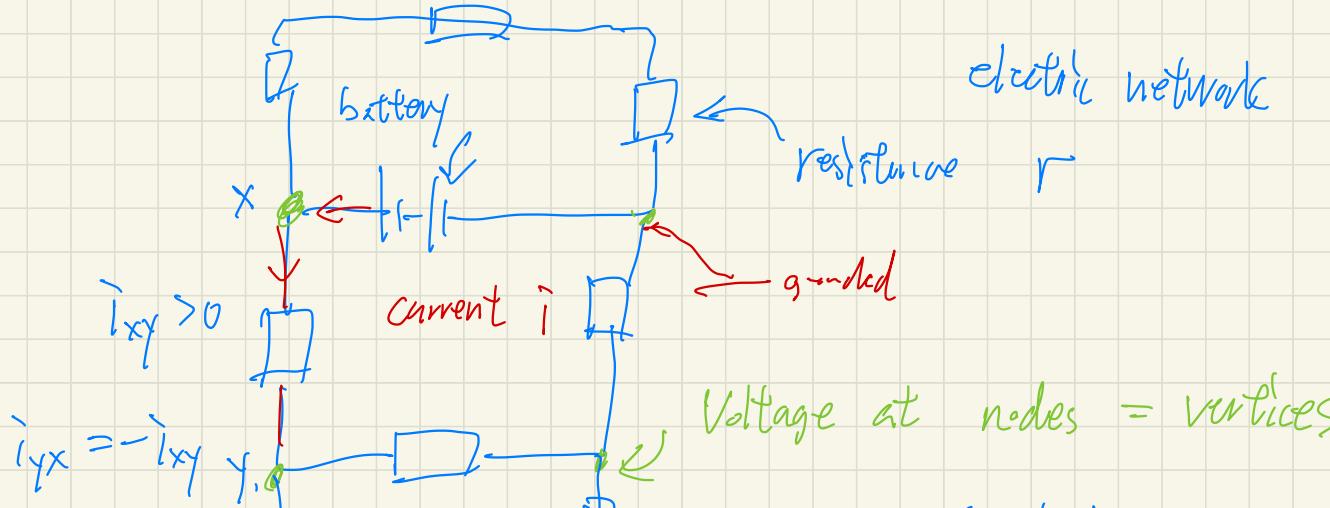
Recall:

$r$  resistance      电 阻  
 $C(x,y) := \frac{1}{r(x,y)}$  conductance

$C(x)$  is called conductance at  $x \in V$ ,

$$D(x) \sim \frac{C(x)}{\sum_{y \in V(x)} C(x,y)}$$

$$\deg(x) < \infty \Rightarrow C(x) < \infty$$



If no external source is present.

Voltage at nodes = vertices.

node:  $\sum \text{outgoing currents} = \sum \text{incoming currents}$   
 (conservation of electrons)

closed loop: sum of voltage around closed loop is zero

if no battery is present.

$$r(x,y) i(x,y) = \text{Voltage drop from } x \text{ to } y.$$

Q1. Represent currents as I-chains.

Dcf. I-chain on  $P$  is real value function over oriented edges s.t.

$$i(x,y) = -i(y,x) \in \mathbb{R} \quad \forall x,y.$$

$i$  is finite if it is nonzero over finitely many edges.

Note:  $i(x,x) = 0$

$\Rightarrow$  self-loops are redundant



(2). Represent external current as 0-chain.

Def. 0-chain is  $j: V \rightarrow \mathbb{R}$ .

$j$  is finitely supported if nonzero over finitely many edges.

Def. Let  $C$  vector space of all 1-chains

s.t.  $\sum_{y \in V(x)} l_i(x,y) < \infty$  if  $x \in V$ .

Def Given 1-chain  $i$ , boundary  $\partial i$  is 0-chain defined by  $(\partial i)_x = \sum_{y \in V(x)} i(x,y)$  i.e.  $\partial i: V \rightarrow \mathbb{R}$  sum of currents at vertex  $i$ .

Def, Given 0-chain  $j$ , boundary  $\partial j \in \mathbb{R}$   
defined by

$$\partial j := \sum_{x \in V} j(x)$$

Ex:  $\partial f_i = 0$  for any 1-chain  $i$ .

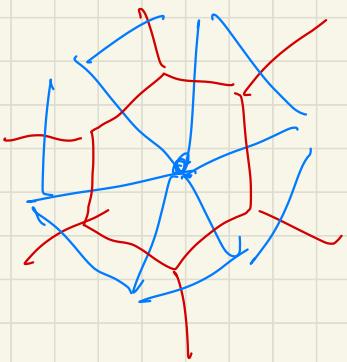
Def 1-chain  $z$  is called a cycle if

$$(\partial z)_x = 0 \quad \forall x \in V.$$

Remark : 1-chain  $\sim$  differential 1-form on  $P^*$

$$i(x,y) = -i(y,x) \sim w(x) = -w(-x) \text{ GLR,}$$

image  $P$  is planar



$$\partial_i = \sum_{y \in V(x)} i(x,y)$$

$$dw = 0 \Leftrightarrow \int\limits_{\gamma} w = 0$$

↑  
exterior derivative

where  $\gamma$  is  
contractible  
closed path.

$P^*$  dual graph

1-cochain  $\sim$  1-form on  $P$

Def (1) A linear functional on vector space of all finite  $l$ -chains is called a  $l$ -cochain.

(2)

-finite 0-chains is ... a cochain.

$$\delta_x(y) = \begin{cases} 1 & \text{if } y=x \\ 0 & \text{if } y \neq x, \end{cases} \quad \text{is 0-chain}$$

$$0\text{-chain: } j = \sum_{x \in V} j(x) \delta_x$$

$\chi_x$  is functional over 0-chain

$$\chi_x(\delta_x) = \begin{cases} 1 & \text{if } x = \tilde{x} \\ 0 & \text{if } x \neq \tilde{x} \end{cases}$$

$$0\text{-cochain: } u(x) = \sum_{x \in V} u(x) \chi_x$$

$$\langle u, j \rangle = \sum_{x \in V} u(x) j(x)$$


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1-chain and 1-cochain:

$$\mathbb{X} = \{x, \bar{x}\}$$

Fix an orientation for edge.

$$X \subset \mathbb{X}$$

$$(x, y) \in X \Rightarrow (y, x) \notin X.$$

Given  $B, D \in \mathbb{X}$ ,  $s_B(D) = \begin{cases} 1 & \text{if } B=D \\ 0 & \text{if } B \neq D \end{cases}$

$$1\text{-chain } i = \sum_{B \in X} i(B) s_B$$

$$\chi_B (\delta_D) = \begin{cases} 1 & \text{if } B=D \\ 0 & \text{if } B \neq D \end{cases}$$

1-cochain  $E = \sum_{BGX} E(B) \chi_B$

$$\langle E, i \rangle = \sum_{BGX} E(B) i(B)$$

Prop : Given 0-cochain  $u$ , coboundary  $\delta^* u$  of  $u$  is the unique 1-cochain s.t,

$$\langle \delta^* u, k \rangle = \langle u, \delta k \rangle$$

for all finite 1-chain  $k$ .

Ex.

$$u: V \rightarrow \mathbb{R},$$

$$\delta^* u(S_{xy}) = u_y - u_x$$

$$\delta^* u(xy)$$

Math message:

Summation at vertices  $\nabla$

boundary

difference along edges  $\delta^*$

co boundary.